

KARYON: Kernel-Based Architecture for safety-critical cOntrol

Cooperation with Disagreement Correction in the
Presence of Communication Failures

KARYON Workshop, Borås, Sweden, Dec 11, 2014



Kernel-Based ARchitecture for safety-critical cONTrol

Introduction

- ▶ Self-driving vehicles plan their trajectories
 - by using **sensory information** about their vicinity
- ▶ Cooperative vehicular systems
 - use V2X for **sharing that sensory information**
- ▶ The limitation is that V2X is **failure** prone
 - and thus we **cannot bound the communication delay**

How can **safety**-critical cooperative systems attain the highest performance in the presence of communication **failures**?

Motivation

- ▶ The cooperative system computes its trajectory based on exchanged information
 - $s = \bigcup_{v \in \text{vehicles}} \{s_v\}$
 - where s_v can be
 - LoS, acceleration, speed, etc
- ▶ Consider two vehicles A and B in platooning
 - Suppose that A does **not** receive s_B but B receives s_A
 - Could A compute its **safe** trajectory according to $\{s_A\}$ and B according to $\{s_A\} \cup \{s_B\}$?
 - This would inevitably lead to **disagreement** about s !

Surprisingly, this could be made safe!

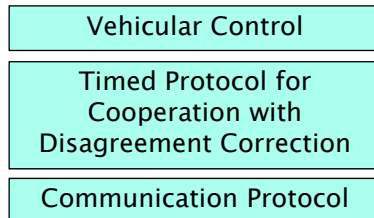
Problem Definition

- ▶ Minimum Longest Uncertainty Period: Is there an upper-bound on the longest period that cooperative systems spend **disagreeing** on s ?
- ▶ This period is longer than **zero**!

We show a solution that uses **1 communication round** in which, in the absence of failures, all vehicles exchange messages.

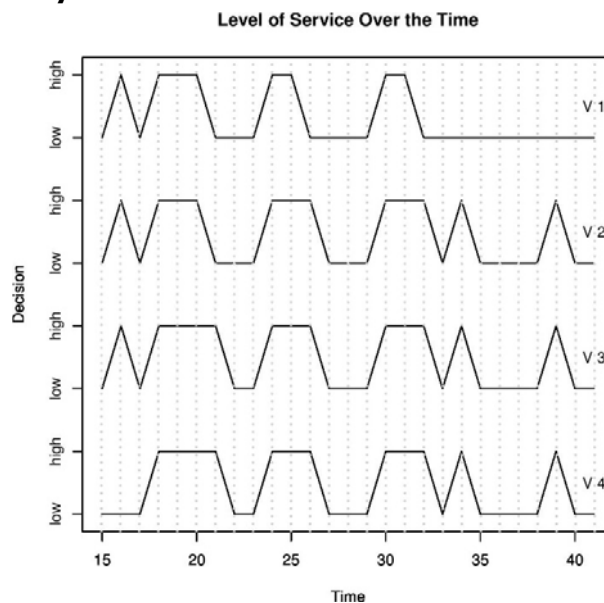


Timed Protocol for Cooperation with Disagreement Correction



Correctness

- ▶ **Theorem 1.** The system disagreement period is bound by one communication round

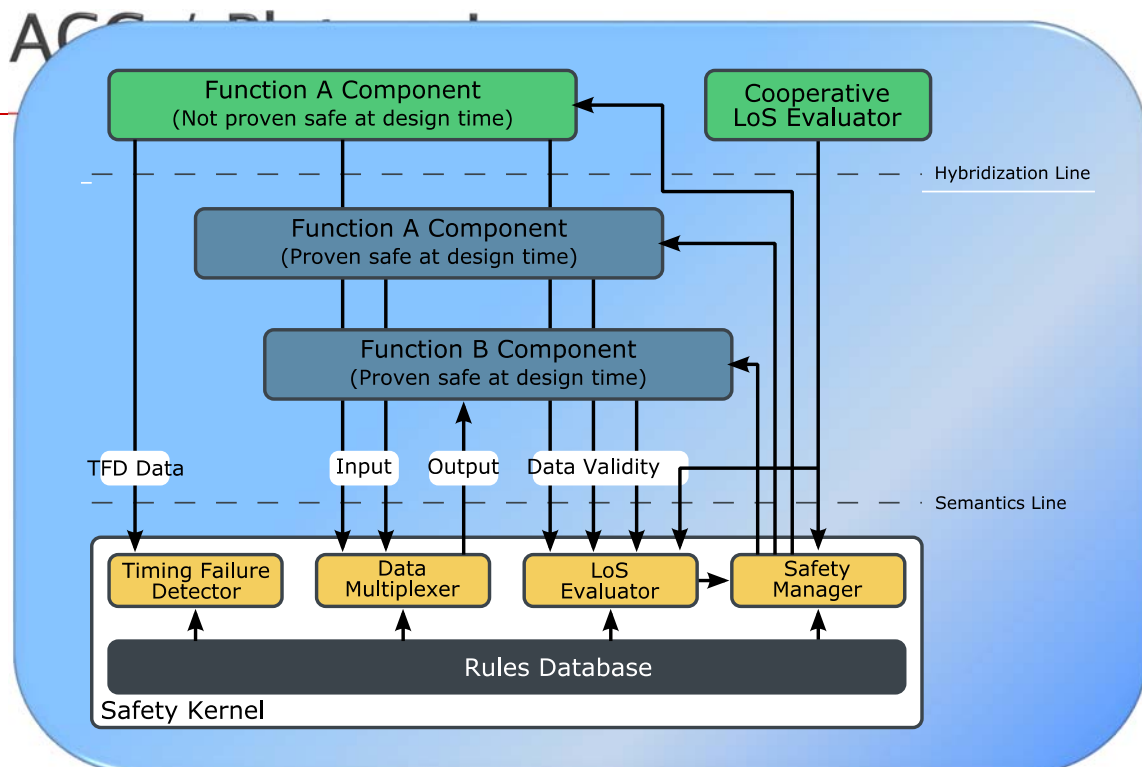
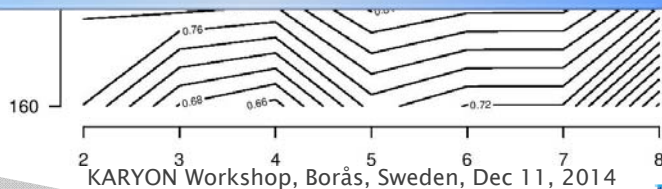


Implementation and Evaluation

- ▶ **Authenticity:** ratio of rounds in which the system **agrees** on s (without using \perp)
 - We use a gossip algorithm with 2 retransmissions

Trade-off between round length and authenticity

- **Shorter** rounds, **shorter** disagreement period but **lower** authenticity
- **Longer** rounds, **longer** disagreement but **higher** authenticity



Conclusion

- We show how to attain the highest performance in the presence of failures
 - Resolves disagreements in at most one round
 - Bounded exposure to risk (that is due to failures)
 - Trade-off: **authenticity** and **disagreement** period
 - Simulations show high **performance** (and still **safe**)
- We demonstrate cooperative applications

Ponce, Schiller, Falcone, “Cooperation with Disagreement Correction in the Presence of Communication Failures,”
17th Inter. IEEE Intelligent Transportation Systems (ITSC’14), CoRR abs/1408.7035